Spectrum Assignment with Non-deterministic Bandwidth of Spectrum Hole in Cognitive Radio Networks

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Abstract. The spectrum allocation for cognitive radio networks (CRNs) has received considerable studies under the assumption that the bandwidth of spectrum holes is static. However, in practice, the bandwidth of spectrum holes is time-varied due to primary user/secondary user (PU/SU) activity and mobility, which result in non-determinacy. This paper studies the spectrum allocation for CRNs with non-deterministic bandwidth of spectrum holes. We present a novel probability density function (PDF) model through order statistic to describe the non-deterministic bandwidth of spectrum holes and provide a bound to approximate it. After that, a statistical spectrum allocation model based on stochastic multiple knapsack problem (MKP) is established for spectrum allocation with non-deterministic bandwidth of spectrum holes. To reduce the computational complexity, we transform this stochastic programming problem into a constant MKP though exploiting the properties of cumulative distribution function (CDF), which can be solved via MTHG algorithm by using auxiliary variable. Simulation results illustrate that the proposed statistical spectrum allocation algorithm can achieve better performances compared to the existing algorithms when the bandwidth of spectrum holes is time-varied.

Keywords—cognitive radio; time-varied; spectrum allocation; non-deterministic bandwidth of spectrum holes; stochastic programming.

1. Introduction

Radio spectrum is becoming one of the most important and scarcest resources in wireless communication system. However, because of the current spectrum allocation policies, the spectrum utilization efficiency in licensed spectrum is very low [1], which generates many non-continuous vacant frequency bands, referred to as spectrum holes [2]. Cognitive radio (CR) which allows secondary users (SUs) to opportunistically utilize the frequency spectrum originally assigned to licensed primary users (PUs) is a promising approach to alleviate spectrum scarcity.

Considerable achievement have been made for spectrum allocation in CRNs by adopting game theory [3]-[4], graph theory [5]-[6] and linear programming [7]-[9], which assume that the parameters of radio environment are static. However, in practice, due to the imperfect spectrum sensing, transmission delay and time-varied spectrum environment etc., it is difficult for the secondary network to have the accurate real-time parameters in CRNs, which will make the parameters non-deterministic. Therefore, in fact, static spectrum allocation algorithm without considering non-deterministic parameters may result in frequent spectrum collision and poor performance. Recently, dynamic resource allocation with non-deterministic parameters in CRNs has received considerable interest from academia, which mainly focus on non-deterministic channel gain and mutual interference [10]-[12]. [10] studies the resource allocation for CRNs under non-deterministic signal-to-interference-plus-noise ratio (SINR) and proposes a power control scheme by using water filling algorithm and stochastic programming. On the other hand, [11] studies the distributed resource allocation problem in CRNs by considering the non-deterministic channel gain and the authors propose a robust distributed power control algorithm by applying second order cone programming (SOCP). [12] proposes a robust distributed uplink power allocation algorithm by using worst case robust optimization method, which consider channel gains from SUs to PUs’ base station, and interference caused by PUs to the SUs’ base station are non-deterministic. In addition, most existing works mainly focus on non-deterministic channel gain and mutual interference. However, due to PU/SU activity and mobility, spectrum environment will change frequently, which result in time-
In this section, the non-deterministic bandwidth of spectrum holes has been formulated in a mathematical way. We consider the bandwidth of spectrum holes as a part of interval between two adjacent order statistics and represent it through a random variable. Then, base on the properties of two adjacent order statistics, the PDF of bandwidth of spectrum holes can be derived.

**Theorem 1:** The PDF of bandwidth of spectrum holes when \( x(i) \) follows a uniform distribution \( U(0,T) \) and \( y(i) \) follows a conditional uniform distribution \( U_j(0, x(i+1) - x(i)) \) is given by

\[
f_x(b) = \frac{\lambda e^{-\lambda b}}{\lambda} \int_0^b f_{x,y}(u,v) du \]

where \( \lambda = \frac{n}{T} \), and the upper and lower bounds for \( f_x(b) \) is

\[
\frac{1}{2} \lambda e^{-2\lambda b} \ln(1 + \frac{2}{\lambda b}) < f_x(b) < \lambda e^{-\lambda b} \ln(1 + \frac{1}{\lambda b})
\]

(This derivation method can also be used when \( x(i) \) and \( y(i) \) follow other distributions)

**Proof:** Let \( z(i) = x(i+1) - x(i) \), where \( z(i) \) represents the interval between two adjacent order statistics. Then the bandwidth of spectrum holes can be achieved by \( b(i) = z(i) - y(i) \). According to [15], the joint probability distribution function of two order statistics can be obtained as

\[
f_{x,y}(u,v) = \frac{n!}{(i-1)!(j-1-i)!(n-j)!} f_x(u) \\
f_x(v) [F_x(u) - F_x(u)^{i-1}] [F_x(v) - F_x(u)]^{j-1} \]

\[
[1 - F_x(v)]^{n-j} - \infty < u < v < \infty
\]

When \( x(i) \) follows uniform distribution, \( f_x(x) = \frac{1}{T} \), \( F_x(x) = \frac{x}{T} \), the probability distribution function \( z \) can be expressed as

\[
f_z(z) = \int_{-\infty}^{\infty} f_{x,y}(u,v) du = \frac{n!}{T (i-1)!(n-i-1)!} \left( 1 - \frac{z}{T} \right)^{i-1} \left( 1 - \frac{z}{T} \right)^{n-i-1} dt
\]

Substituting \( B(\alpha, \beta) = \frac{\Gamma(\alpha)\Gamma(\beta)}{\Gamma(\alpha + \beta)} \) in (4) yields

\[
f_z(z) = \frac{n}{T} \left( 1 - \frac{z}{T} \right)^{i-1} \approx \frac{n}{T} \left( 1 - \frac{z}{T} \right)^{i-1} \approx \frac{n}{T} \left( 1 - \frac{z}{T} \right)^{i-1} \approx \frac{n}{T} \left( 1 - \frac{z}{T} \right)^{i-1}
\]

3. **Statistical Model for Bandwidth of Spectrum Holes**
Since \( y(i) \) follows a conditional uniform distribution, uniformly distributed in \((0, z)\), according to Bayes' theorem, the joint probability distribution function \( f_{y,z}(y,z) \) can be obtained as

\[
f_{y,z}(y,z) = \frac{1}{z} e^{-\frac{n}{T}}
\]

(6)

The probability distribution function of \( b \) can be given as

\[
f_b(b) = \int f_{y,z}(y,b)dy = \frac{e^{-\frac{n}{T}b}}{T} du
\]

(7)

Substituting exponential integral function,

\[
E_t(ax) = \int_{x}^{\infty} e^{-ax} dx
\]

in (7), result in

\[
f_b(b) = \lambda [E_1(\lambda b) - E_1(n)]
\]

(8)

Let \( \lambda = \frac{n}{T} \), since the value of \( E_1(n) \) is minimal, (8) can be approximated as \( \lambda E_1(\lambda b) \). Based on the properties of exponential integral function[16], we obtain the upper and lower bounds for \( f_b(b) \) is given as

\[
\frac{1}{2}\lambda e^{-\lambda b} \ln(1 + \frac{2}{\lambda b}) < f_b(b) < \lambda e^{-\lambda b} \ln(1 + \frac{1}{\lambda b})
\]

(9)

4. Statistical Spectrum Allocation Algorithm Based on Stochastic MKP

4.1 Spectrum Allocation Model Based on Stochastic MKP

We assume that the CRNs consist of \( m \) SUs and \( n \) spectrum holes, where SUs’ data rate requirements and the bandwidth of spectrum holes are indicated as \( R = [r_1, r_2, \ldots, r_m] \) and \( C = [c_1, c_2, \ldots, c_n] \) respectively. Among this, each SU \( i \) (\( i = 1, 2, 3, \ldots, m \)) can utilize any spectrum hole \( j \) (\( j = 1, 2, 3, \ldots, n \)), which means the non-continuous spectrum allocation problem can be transformed into a kind of MKP to achieve maximum transmission rate.

\[
\text{maximize} \quad z = \sum_{i=1}^{m} \sum_{j=1}^{n} r_i x_{i,j}
\]

(10)

\[
\text{s.t.} \quad \sum_{j=1}^{n} b_j x_{i,j} \leq c_i \quad j \in N = \{1, \ldots, n\}
\]

(11)

\[
\sum_{j=1}^{n} x_{i,j} \leq 1 \quad i \in M = \{1, \ldots, m\}
\]

(12)

\[
x_{i,j} = 0 \text{ or } 1 \quad i \in M \quad j \in N
\]

(13)

Let \( b \) be the necessary bandwidth that satisfied data rate requirement \( r \) and \( x_{i,j} = 1 \) represent SU \( i \) can allocated into spectrum hole \( j \). Moreover, by adopting AMC scheme, SUs with different positions can achieve different data rate in same channel due to environment differences. According to G.J. Foschini’s inference[17], the bandwidth needs for SU \( i \) on spectrum hole \( j \) can be given as

\[
b_{i,j} = \frac{r_i}{\log_2(1 + \frac{\gamma_{i,j}}{P_b}) - \frac{2}{3} \ln \frac{\gamma_{i,j}}{P_b}}
\]

(14)

Where \( P_b \) denotes maximum tolerable error rate and \( \gamma_{i,j} \) indicates the SINR when SU \( i \) transmits on spectrum hole \( j \).

Substituting (14) into (11), we obtain a special MKP model, where each SU requires different bandwidth when allocated into different spectrum holes. Then (11) can be expressed as

\[
\sum_{j=1}^{n} \frac{b_{i,j}}{ca_{i,j}} x_{i,j} \leq \frac{c_i}{ca_{\text{sum}}} \quad j \in N = \{1, \ldots, n\}
\]

(15)

Where \( ca_{i,j} \) is the subcarrier bandwidth, \( ch_a \) denotes the channel bandwidth, \( ca_{\text{sum}} \) indicates the number of subcarriers in each channel.

To consider non-deterministic bandwidth of spectrum holes, we take previous statistical model into this MKP model. Let \( c_{i,j} \) be the available bandwidth of spectrum hole \( j \) detected by SUs’ base station at time \( t \) and \( C_j = \{c_{i,j-k+1}, c_{i,j-k+2}, \ldots, c_{i,j}\} \) represents the record of bandwidth of spectrum hole \( j \) in recent \( k \) times spectrum sensing. We can determine the parameter \( \lambda \) of the statistical model by applying point estimation method. The expectation of previous PDF can be written as

\[
E(b) = \int_{0}^{n} E_t(\lambda b) db = -\lambda \cdot \frac{\Gamma(2)}{2 \lambda^2} = \frac{1}{2 \lambda}
\]

(16)

Hence, based on the method of moment[15], the PDF parameter \( \lambda \) of bandwidth of spectrum hole \( j \) can express as \( 1/2E(C_j) \).

We assume \( c_i \) is a random variable in (15) and follows the PDF (1), which transmits the MKP spectrum allocation model above into a stochastic programming model. The unknown random variable introduces uncertainty to (15) and makes the constraint quit complicated. Thus, to reduce the complexity, we tackle this issue by using chance-constraint programming method which finds the optimum solution under the condition that the constraints is feasible with probability greater than the threshold \( \alpha \). The constraint (15) above can be rewritten as

\[
\text{s.t.} \quad \Pr \left\{ \frac{\sum_{i=1}^{m} b_{i,j} x_{i,j}}{ca_{i,j}} \leq \frac{c_i}{ca_{\text{sum}}} \right\} \geq \alpha \quad j \in N
\]

(17)
This indicates the constraint must be feasible with the probability greater than \( \alpha_j \), and this model maximizes the transmission rate by satisfying the constraint with high probability.

4.2 Simplification and solution

To reduce the computational complexity, we simplify the stochastic programming model above by the usage of the properties of CDF. For (17), since \( c_j \) follows the CDF \( F_j(c) \), there must exist a \( K_j \) for each threshold \( \alpha_j \) (0 \( \leq \alpha_j \leq 1 \)) satisfies \( \Pr\{K_j \leq c_j\} = \alpha_j \). Therefore, (17) will be satisfied only if

\[
\sum_{i=1}^{n} \frac{b_{i,j} \cdot c_{i,j}}{ca_i \cdot ca_{num}} \cdot x_{i,j} \leq K_j
\]

Furthermore, as \( \Pr\{K_j \leq c_j\} = 1 - F_j(K_j) \), \( K_j \) can be represented by \( F_j^{-1}(1-\alpha_j) \). We can transmit (17) as

\[
\sum_{i=1}^{m} \frac{b_{i,j} \cdot c_{i,j}}{ca_i \cdot ca_{num}} \cdot x_{i,j} \leq F_j^{-1}(1-\alpha_j) \quad j \in N = \{1,...,n\}
\]

Since the model proposed is a NP-hard problem, the existing algorithm obtains the solution by exhaustive search which will encounter complex computations when \( n \) and \( m \) have large values. Hence, we convert the model into a Generalized Assignment Problem (GAP) by using auxiliary variable and solve it via MTHG algorithm [17] which will get the suboptimal solution. Let

\[
\frac{b_{i,j} \cdot c_{i,j}}{ca_i \cdot ca_{num}} = w_{i,j}, \quad \tilde{c}_j = F_j^{-1}(1-\alpha_j), \quad p_{i,j} = r, \quad j \in N
\]

and adding auxiliary variables \( p_{i,j} = 0 \), \( \tilde{c}_{i+1} = m \) and \( w_{i,1} = 1, i \in M \). Then, the model can be rewritten as

\[
\text{maximize } z = \sum_{i=1}^{n} \sum_{j=1}^{m} p_{i,j} \cdot x_{i,j}
\]

s.t.

\[
\sum_{j=1}^{m} w_{i,j} \cdot x_{i,j} \leq \tilde{c}_j \quad j \in \tilde{N} = \{1,...,n,n+1\}
\]

\[
\sum_{j=1}^{m} x_{i,j} = 1 \quad i \in M = \{1,...,m\}
\]

\[
x_{i,j} = 0 \text{ or } 1 \quad i \in M \quad j \in \tilde{N}
\]

This GAP model means each SU must be allocated in a unique \( \tilde{c}_j \) and \( \tilde{c}_{i+1} \) contains the SUs which have not be allocated due to lack of spectrum resource. The GAP model can be solved by using MTHG algorithm, where iteratively consider all the unassigned SUs, and assign the SU \( i' \) having the maximum difference between the largest and second largest \( p_{i,j} \) to the spectrum hole \( j' \) for which \( p_{i,j} \) is a maximum (\( p_{i,j'} = \max\{p_{i,j'}: j \in \tilde{N}\}\)).

5. Simulation Results

In this section, the performance of proposed statistical spectrum allocation algorithm has been studied through numerical simulations. We assume that all SUs located around SUs’ base station with the distance randomly ranged from 50 to 2000 m and the simulation parameters are set as Table I. In order to simulate the time-varied bandwidth of spectrum holes caused by PU/SU activity and mobility, we randomly generate spectrum holes with the bandwidth changed through change rate \( p \) (the probability of bandwidth of spectrum hole changed next moment) over the frequency band \( W \). Then, we compare the performance of proposed statistical MKP spectrum allocation algorithm with static MKP spectrum allocation algorithm in [8].

<table>
<thead>
<tr>
<th>Parameters</th>
<th>Values</th>
</tr>
</thead>
<tbody>
<tr>
<td>SU output power ( P_{\text{out}} )</td>
<td>45dBm</td>
</tr>
<tr>
<td>Maximum tolerable error rate ( P_{\text{e}} )</td>
<td>10^-4</td>
</tr>
<tr>
<td>SU number ( SMUN )</td>
<td>30</td>
</tr>
<tr>
<td>Spectrum hole number ( SH )</td>
<td>3~6</td>
</tr>
<tr>
<td>Noise power ( N_{\text{B}} )</td>
<td>-111dBm</td>
</tr>
<tr>
<td>Chance programming threshold ( \alpha )</td>
<td>0.5</td>
</tr>
<tr>
<td>Bandwidth of spectrum hole change rate ( p )</td>
<td>0.1~1</td>
</tr>
<tr>
<td>Total frequency band ( W )</td>
<td>10MHZ</td>
</tr>
</tbody>
</table>

Fig. 2 illustrates the CDF of SU spectrum efficiency using the two algorithms with different situations (\( p = 0.5 \) and \( p = 1 \)). As shown in Fig. 2, statistical MKP algorithm and static MKP algorithm occupy 61% and 52% respectively when the spectrum efficiency is large than 1.8, which means statistical MKP algorithm has larger percentage of high spectrum efficiency than the static MKP algorithm. This case is because statistical MKP algorithm which takes the statistical properties of bandwidth of spectrum holes into consideration can achieve dynamic optimal and reduce bandwidth collisions caused by spectrum holes changes. Moreover, as \( p \) increased to 1, statistical MKP algorithm exceeds static MKP algorithm more, which means statistical MKP algorithm can better adapt to the scenario with frequency time-varied bandwidth of spectrum holes and achieve high spectrum efficiency.
In Fig.3, we evaluate the network utility with different bandwidth change rate $p$. The network utility is defined as [18]

$$U = \sum_{k} \ln R(k)$$

(24)

Where $R(k)$ is the throughput for SU $k$. In the Fig. 3, as $p$ increased, the network utility is reduced due to bandwidth collision. The static MKP algorithm suffers lower network utility than statistical MKP algorithm, because it only considers current static optimal and use constant allocation parameters, which may lead to more bandwidth collision. In addition, when the number of spectrum holes increase, the network utility of statistical MKP algorithm increase more than the static MKP algorithm, which means statistical MKP algorithm can obtain good performance under the condition of non-continuous spectrum.

Figure 3. Network utility with different bandwidth of spectrum holes change rate $p$

Fig. 4 shows the average spectrum collision rate with different bandwidth change rate $p$. From Fig. 4, we can see that as $p$ increase, the spectrum collision rate increased dramatically and statistical MKP algorithm achieve lower collision rate than the static MKP algorithm. It is because statistical MKP algorithm can relax the bandwidth collision through assigning spectrum resource with considering the statistical properties of spectrum holes. The result shows that statistical MKP algorithm can effectively reduce the spectrum collision when spectrum holes changes.

Figure 4. Average spectrum collision rate with different bandwidth of spectrum holes change rate $p$

6. Conclusion

In this paper, we study the problem of non-continuous spectrum allocation in CRNs where the bandwidth of spectrum holes is non-deterministic due to PU/SU activity and mobility. We present a novel PDF model through order statistic to describe the non-deterministic bandwidth of spectrum holes and provide a bound to approximate it. After that, a statistical spectrum allocation model based on stochastic MKP is established for spectrum allocation with non-deterministic bandwidth of spectrum holes. To reduce the computational complexity, we transform this stochastic programming problem into a constant MKP through exploiting the properties of CDF, which can be solved via MTHG algorithm by using auxiliary variable. Simulation results verify that the proposed statistical spectrum allocation algorithm can achieve better performances compared to the existing algorithms when the bandwidth of spectrum holes is time-varied.

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